

# Cryptosystem of Chua and Ling

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The authors show that the cubic curve cryptosystem proposed by Chua and Ling can easily be reduced to the cryptosystem of Rabin-Williams.

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*Introduction:* At Eurocrypt '96, Meyer and Müller [1] presented a new Rabin-type scheme based on elliptic curves. In [2], this system was reduced to the system of Rabin-Williams [3, 4]. Using the same technique, we show that the system of Chua and Ling [5] can also be reduced.

*Chua-Ling's cryptosystem:* This cryptosystem is based on a singular cubic curve of the form

$$C_n(b) : y^2 \equiv x^3 + bx^2 \pmod{n}.$$

To setup the system, each user chooses two large primes  $p$  and  $q$  both congruent to 11 modulo 12. Then, he publishes the value of  $n = pq$ .

Suppose Alice wants to send a message  $m$  to Bob. Then she chooses  $\lambda \in \mathbb{Z}/n\mathbb{Z} - \{0, \pm 1\}$  and sets  $\mathbf{P} = (m^2, \lambda m^3)$ . Next, she computes  $c = \lambda^3 \pmod{n}$  and  $b = (\lambda^2 - 1)m^2 \pmod{n}$ , and sends the ciphertext consisting of  $c, b, x_Q = x([2]\mathbf{P}), t = \left(\frac{y([2]\mathbf{P})}{n}\right)$  and  $u = \text{lsb}(y([2]\mathbf{P}))$ .

To recover the plaintext  $m$ , Bob computes the unique  $y_Q$  satisfying  $y_Q^2 \equiv x_Q^3 + bx_Q^2 \pmod{n}$  with Jacobi symbol  $t$  and  $\text{lsb } u$ . He sets  $\mathbf{Q} = (x_Q, y_Q)$ . Letting  $\mathbf{Q}_p = \mathbf{Q} \pmod{p}$  and  $\mathbf{Q}_q = \mathbf{Q} \pmod{q}$ , he computes  $\mathbf{P}_{i,p} = (x_{i,p}, y_{i,p})$  ( $i = 1, 2$ ) such that  $[2]\mathbf{P}_{i,p} = \mathbf{Q}_p$  on  $C_p(b)$  and similarly  $\mathbf{P}_{i,q}$ . Next, he computes  $I_p = \{i : c^2 \equiv y_{i,p}^6 x_{i,p}^{-9} \pmod{p}\}$ . He does the same for the prime  $q$ . Finally, he computes  $m_p = y_{i,p}^3 x_{i,p}^{-4} c^{-1} \pmod{p}$  ( $i \in I_p$ ) and  $m_q$ . So, Bob obtains  $m$  using the Chinese Remainder Theorem such that  $m \equiv m_p \pmod{p}$  and  $m \equiv m_q \pmod{q}$ .

*Reduction to Rabin-Williams:* Since  $\mathbf{P} = (m^2, \lambda m^3) \in C_n(b)$ , it follows that:

$$x_Q = x([2]\mathbf{P}) \equiv \frac{(3m^2 + 2b)^2}{4\lambda^2 m^2} - 2m^2 - b \pmod{n}, \quad (1)$$

and

$$\lambda^2 m^2 \equiv m^2 + b \pmod{n}. \quad (2)$$

From eqns. 1 and 2, we construct the polynomials  $\mathcal{P}_1, \mathcal{P}_2 \in (\mathbb{Z}/n\mathbb{Z})[X]$  given by

$$\begin{aligned} \mathcal{P}_1(X) &= 4(x_Q + 2X + b)(X + b) - (3X + 2b)^2, \\ \mathcal{P}_2(X) &= c^2 X^3 - (X + b)^3. \end{aligned}$$

Since  $m^2$  is a root of  $\mathcal{P}_1$  and  $\mathcal{P}_2$ ,  $m^2$  will be a root of  $\mathcal{R} = \text{gcd}(\mathcal{P}_1, \mathcal{P}_2)$ . The polynomial  $\mathcal{R}$  is very likely to be of degree 1 [6]. Solving this polynomial in  $X$  gives the value of  $m^2$ .

*Conclusion:* We have shown that we can easily recover the value of  $m^2$  from the ciphertext corresponding to a plaintext  $m$ . Therefore, the Chua-Ling scheme is reduced to the Rabin-Williams cryptosystem.

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